

Revisiting Optimal Rank Aggregation: A Dynamic Programming Approach



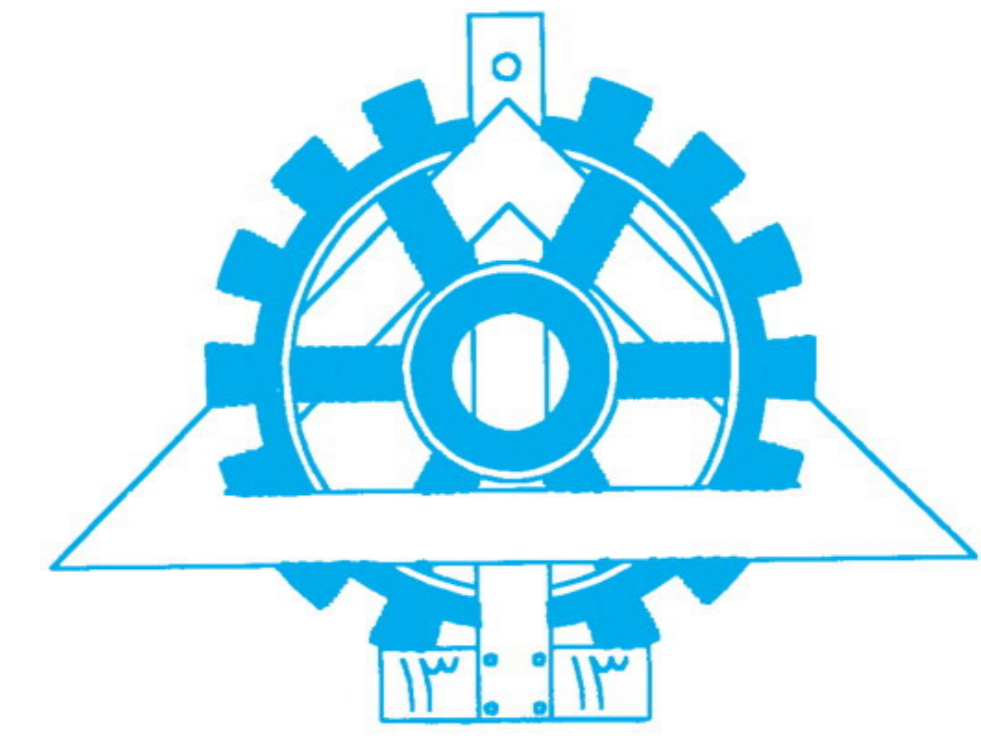
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Problem

One of the most important applications of rank aggregation [3], that is merging multiple ranked lists of objects, is in the field of information retrieval (IR). Typically, rank aggregation in IR is conducted for two different purposes:

1. Aggregating ranked lists of a document set ranked by different ranking methods.
2. Aggregating ranked lists of different document sets ranked by the same ranking method. \triangleright Our goal!

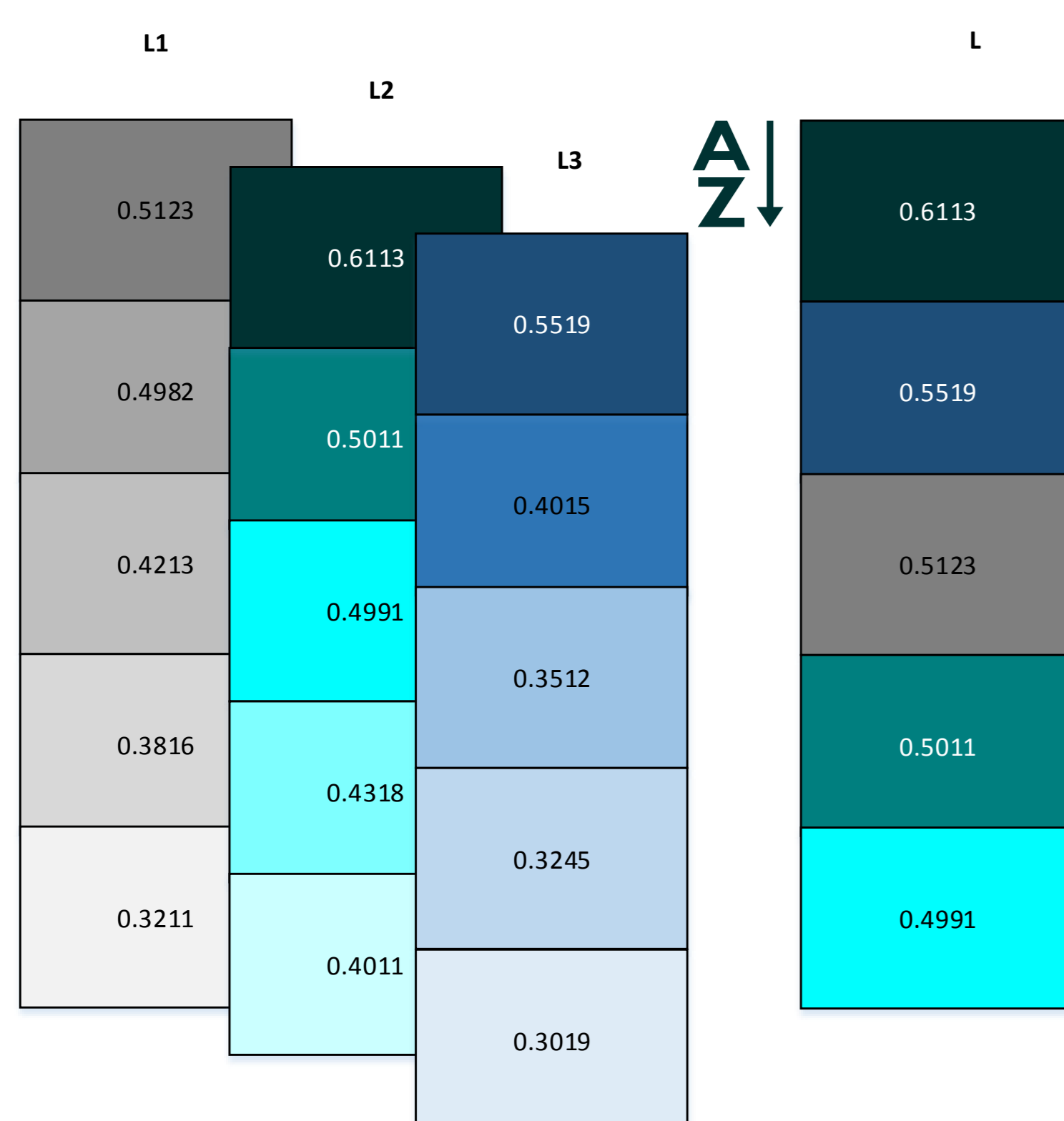


FIGURE 1: Rank aggregation problem: aggregating ranked lists of different document sets ranked by the same ranking method

Optimal Rank Aggregation: a rank aggregation problem in which, with the presence of relevancy information, items from different ranked lists are aggregated into a single final list so that:

- The final ranking has the best possible achievable performance;
- Local ordering of the initial lists are preserved in the final list.

Chen's Algorithm: Bucketing Strategy

A **bucket** is denoted in the form of $(n^-, n^+)\{d_i, d_{i+1}, \dots, d_j\}$, in which n^- and n^+ respectively determine the number of **irrelevant** and **relevant** documents in the bucket and $\{d_i, d_{i+1}, \dots, d_j\}$ denotes the documents ranked in positions i to j . For example $(3, 1)\{d_3, d_4, d_5, d_6\}$ has 3 irrelevant documents, $d_3, d_4,$ and d_5 and one relevant document d_6 .

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1: procedure CHENAGGREGATOR(Initial Ranked Lists)
2:   for each List  $l_i$  in the Initial Ranked Lists do
3:      $b_i \leftarrow$  BucketRepresentation( $l_i$ );
4:   end for
5:   MergedList  $\leftarrow$   $\emptyset$ 
6:   ActiveSet  $\leftarrow$   $\{b_i^1 \mid 1 \leq i \leq l \wedge n_{b_i^1}^+ \geq 1\}$ 
7:   repeat
8:     ActiveSet  $\leftarrow$  Sort(ActiveSet)
9:      $b_k^j \leftarrow$  ActiveSet[1]
10:    MergedList.append( $b_k^j$ )
11:    ActiveSet.remove( $b_k^j$ )
12:    if  $n_{b_k^j}^+ \geq 1$  then
13:      ActiveSet.add( $b_k^{j+1}$ )
14:    end if
15:  until isEmpty(ActiveSet)
16:  return MergedList
17: end procedure

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FIGURE 2: Chen's algorithm [1, 2] for optimal rank aggregation

In this example, based on Chen's algorithm, the initial *active set* will be: $\{(0, 1)\{A_1\}, (2, 1)\{B_1, B_2, B_3\}, (1, 3)\{C_1, C_2, C_3, C_4\}\}$. Continuing the algorithm, it will produce the following final result: $\{\{A_1\}, \{C_1, C_2, C_3, C_4\}, \{A_2, A_3\}, \{B_1, B_2, B_3\}, \{A_4\}, \{B_4\}\}$.

TABLE 1: The results of three different rankings. The relevant and irrelevant documents are marked by '+' and '-' respectively.

Rank	A	B	C
1	A_1^+	B_1^-	C_1^-
2	A_2^-	B_2^-	C_2^+
3	A_3^+	B_3^+	C_3^+
4	A_4^-	B_4^-	C_4^+

TABLE 2: Buckets of documents for the ranked lists.

Set	A	B	C
1	$(0, 1)\{A_1\}$	$(2, 1)\{B_1, B_2, B_3\}$	$(1, 3)\{C_1, C_2, C_3, C_4\}$
2	$(1, 1)\{A_2, A_3\}$	$(1, 0)\{B_4\}$	
3	$(1, 0)\{A_4\}$		

Counter-example

- The greedy approach does not necessarily lead to an optimal solution.
- In this example, the average precision of Chen's algorithm output is 0.5289, while that of the optimal aggregated list is 0.6198.

TABLE 3: The counterexample that shows that Chen's algorithm does not necessarily generate the optimal aggregated list.

Rank	A	B	Chen	DP
1	A_1^-	B_1^-	A_1^-	B_1^-
2	A_2^+	B_2^-	A_2^+	B_2^-
3	A_3^-	B_3^+	A_3^-	B_3^+
4	A_4^+	B_4^+	A_4^+	B_4^+
5	A_5^-	B_5^+	A_5^-	B_5^+
6	A_6^+	B_6^+	A_6^+	B_6^+
7	A_7^-	B_7^+	A_7^-	B_7^+
8	A_8^+	B_8^+	A_8^+	B_8^+
9			B_1^-	A_1^-
10			B_2^-	A_2^+
11			B_3^+	A_3^-
12			B_4^+	A_4^+
13			B_5^+	A_5^-
14			B_6^+	A_6^+
15			B_7^+	A_7^-
16			B_8^+	A_8^+
MAP			0.5289	0.6198

Proposed DP Algorithm

We introduce an m -th order tensor M of size $(|b_1| + 1) \times (|b_2| + 1) \times \dots \times (|b_m| + 1)$, in which value of an entry $M[i_1, i_2, \dots, i_m]$ shows the maximum AP obtainable by aggregating top i_1, i_2, \dots, i_m buckets of the ranked lists $1, 2, \dots, m$, respectively.

$$M[i_1, i_2, \dots, i_m] = \max_k (M[i_1, i_2, \dots, i_k - 1, \dots, i_m] + \text{AP-GAIN}(k, i_1, i_2, \dots, i_m)). \quad (1)$$

AP-GAIN(k, i_1, i_2, \dots, i_m) in Eq. 1 is the AP increase after appending b_k^i at position p (i.e. at the end of the list):

$$p = \sum_{1 \leq t \leq m} \sum_{\substack{1 \leq r \leq i_t \\ \neg(t=k \wedge r=i_k)}} n_{b_r^t}. \quad (2)$$

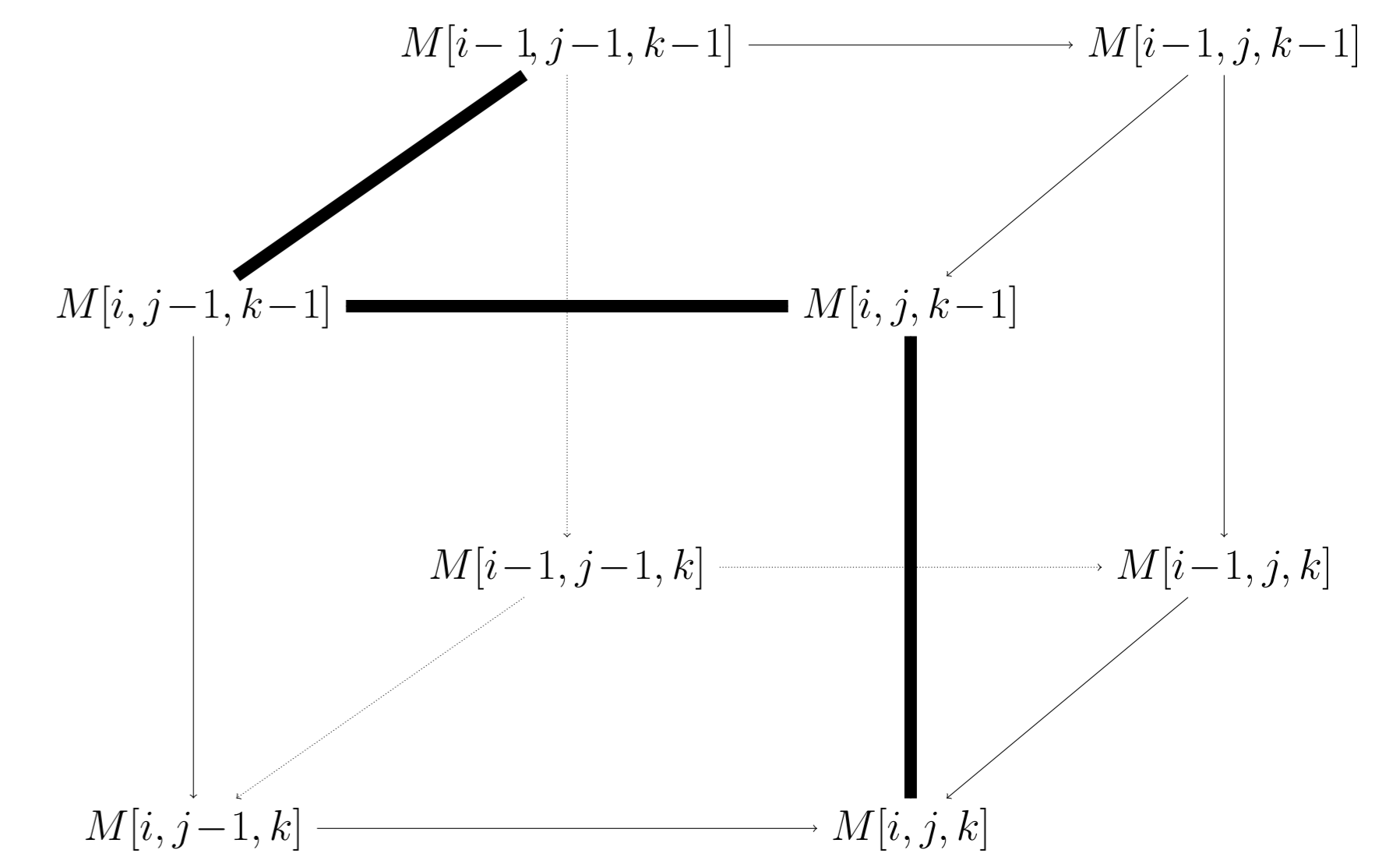


FIGURE 3: A graph of MAP belonging to bucket aggregation of 3 ranked lists. $M[i, j, k]$ shows the maximum AP obtainable by aggregating top $i, j,$ and k buckets the ranked lists 1, 2 and 3 respectively.

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1: procedure DPAGG(Bucket Representation)
2:    $M \leftarrow [0]_{|b_1| \times \dots \times |b_m|}$ 
3:   for  $i_1 \leftarrow 1$  to  $|b_1|$  do
4:     ...
5:     for  $i_m \leftarrow 1$  to  $|b_m|$  do
6:        $v \leftarrow M[i_1 - 1, \dots, i_m] + \text{AP-GAIN}(1, i_1, \dots, i_m)$ 
7:        $k \leftarrow 1$ 
8:       for  $k' \leftarrow 2$  to  $m$  do
9:          $\Delta A_{k'} \leftarrow \text{AP-GAIN}(k', i_1, \dots, i_{k'} - 1, \dots, i_m)$ 
10:         $v' \leftarrow M[i_1, \dots, i_{k'} - 1, \dots, i_m] + \Delta A_{k'}$ 
11:        if  $v' \geq v$  then
12:           $v \leftarrow v'$ 
13:           $k \leftarrow k'$ 
14:        end if
15:      end for
16:       $M[i_1, \dots, i_k - 1, \dots, i_m] \leftarrow v$ 
17:    end for
18:  end for
19: return M

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FIGURE 4: DP algorithm for optimal rank aggregation

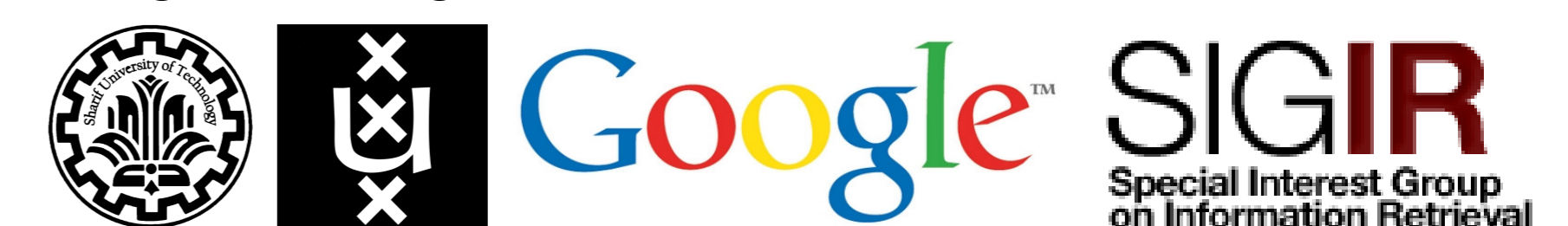
By storing index k maximizing Eq. 1 for each entry $M[i_1, i_2, \dots, i_m]$, the optimal solution could be easily generated by starting from entry $M[|b_1|, |b_2|, \dots, |b_m|]$ and iteratively inserting the bucket indicated by the stored index, at the beginning of the output.

Time complexity of:

- DP algorithm: $O(m \times |b_1| \times |b_2| \times \dots \times |b_m| + n \times t)$.
- Brute-force algorithm: $O\left(\binom{|b_1| + |b_2| + \dots + |b_m|}{|b_1|, |b_2|, \dots, |b_m|}\right)$.

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